

Isolate Vertex Resolving Partition in a Graph

S. Hemalathaa¹, A. Subramanian², P. Aristotle³, V.Swamianthan⁴

¹(Department of Mathematics, The M.D.T. Hindu College, Thirunelveli- 627010, Tamilnadu, India)

²(Department of Mathematics, The M.D.T. Hindu College, Thirunelveli- 627010, Tamilnadu, India)

³(Department of Mathematics, Raja Doraisingam Government Arts College, Sivagangai – 630561, Tamilnadu, India)

⁴(Ramanujan Research Centre in Mathematics, Saraswathi Narayan College, Madurai-625022, Tamilnadu, India)

Abstract: Let $G = (V, E)$ be a simple connected graph. A partition $\pi = \{V_1, V_2, V_3, \dots, V_k\}$ is called a resolving partition of G if for any $u \in V(G)$, the code of u with respect to π (denoted by $c_\pi(u)$) namely $(d(u, V_1), d(u, V_2), \dots, d(u, V_k))$ is distinct for different $u \in V(G)$ where $d(u, V_i) = \min\{d(u, x) / x \in V_i\}$. The minimum cardinality of a resolving partition of a graph G is called the partition dimension of G and is denoted by $pd(G)$. Several types of resolving partition have been considered like connected resolving partition, metric chromatic number of a graph (that is, independent resolving partition), equivalence resolving partition etc. In this paper, a new type of resolving partition is considered and in this partition, each element of the partition contains an isolate vertex in the subgraph induced by that element. This is a generalization of an independent resolving partition. This resolving partition is called an isolate vertex resolving partition in a graph. The minimum cardinality of an isolate vertex resolving partition is denoted by $pd_{is}(G)$. This parameter for some well known graphs is found. Graphs for which $pd_{is}(G) = 2$ or $pd_{is}(G) = n$ are characterized.

Keywords: Central vertex, Isolate vertex partition dimension, Isolate Vertex resolving partition, Partition Dimension, Resolving partition.

1. Introduction

In this paper, G is a simple, finite, connected and undirected graph.

Definition 1.1. [1] Let $G = (V, E)$ be a simple, finite, connected and undirected graph. A partition $\pi = \{V_1, V_2, V_3, \dots, V_k\}$ of $V(G)$ is called a resolving partition of G if the code $c_\pi(u) = (d(u, V_1), d(u, V_2), \dots, d(u, V_k))$ is distinct for different $u \in V(G)$ where $d(u, V_i) = \min\{d(u, x) / x \in V_i\}$. The minimum cardinality of a resolving partition of a graph G is called the partition dimension of G and is denoted by $pd(G)$.

Definition 1.2. Let $G = (V, E)$ be a simple connected graph. Let $\pi = \{V_1, V_2, V_3, \dots, V_k\}$ be a partition of $V(G)$. If each $\langle V_i \rangle$ contains an isolate and if π is a resolving partition, then π is called an isolate vertex resolving partition. The trivial partition namely $\pi = \{\{u_1\}, \{u_2\}, \dots, \{u_n\}\}$ where $V(G) = \{u_1, u_2, \dots, u_n\}$ is an isolate vertex resolving partition.

The minimum cardinality of an isolate vertex resolving partition is called the isolate vertex partition dimension of G and is denoted by $pd_{is}(G)$.

Remark 1.3. Every independent resolving partition is an isolate vertex resolving partition. Therefore $pd_{is}(G) \leq ipd(G) \leq pd(G)$.

2. $pd_{is}(G)$ FOR WELL KNOWN GRAPHS

1. $pd_{is}(K_n) = n$
2. $pd_{is}(K_{1,n}) = n + 1$
3. $pd_{is}(K_{m,n}) = m + n$
4. $pd_{is}(K_{a_1, a_2, \dots, a_n}) = a_1 + a_2 + \dots + a_n$
5. $pd_{is}(K_m(a_1, a_2, \dots, a_m)) = \max\{a_1, a_2, \dots, a_m\}$

Theorem 2.1. $pd_{is}(W_n) \geq 5$ if $n \geq 14$.

Proof: Suppose $n \geq 14$.

Let $\pi = \{\{u\}, V_1, V_2, V_3\}$ be an isolate vertex partition, where u is the central vertex. Every vertex V_i ($1 \leq i \leq 3$) has code 1 with respect to $\{u\}$ and 0 with respect to V_i . Since there are only two more partitions and since the codes with respect to these two partitions can be either (1,1) or (1,2) or (2,1) or (2,2), $|V_i| \leq 4$ for all i . That is

$|V_1 \cup V_2 \cup V_3| \leq 12$. Since there are 13 vertices in $V_1 \cup V_2 \cup V_3$, we get a contradiction. Therefore $pd_{is}(W_n) \geq 5$ if $n \geq 14$.

Remark 2.2.

1. $pd_{is}(W_n) \geq 6$ if $n \geq 34$.
2. $pd_{is}(W_n) \geq 7$ if $n \geq 82$.
3. $pd_{is}(W_n) \geq k + 1$ if $n \geq (2^{k-2} + 1) + (k - 2) 2^{k-2} + 1$.

Theorem 2.3. Let G be a connected graph of order n . Then $pd_{is}(G) = 2$ if and only if $G = K_2$.

Proof: Let $\pi = \{V_1, V_2\}$ be a minimum isolate vertex resolving partition.

When $n = 2$, V_1 and V_2 are singletons and $G = K_2$. When $n = 3$, π is not an isolate vertex resolving partition.

Suppose $n = 4$. Then $G = K_4, C_4, K_{1,3}, K_4 - e, P_4, K_{1,3} + e$ and $pd_{is}(G) = 4$. Let $n \geq 5$. Let $\Pi = \{V_1, V_2\}$ be a minimum isolate vertex resolving partition. V_1 contains an isolate vertex say z . Then $d(z, V_2) = 1$. Suppose V_1 contains three vertices say z, x and y . $c_{\pi}(z) = (0, 1)$. Therefore $c_{\pi}(x) = (0, t)$, where $t \neq 1$ and $c_{\pi}(y) = (0, t)$, where $t \neq 1, t \neq t$. Therefore one of $t, t \geq 3$. Let $t \geq 3$. Then there exists a shortest path

$y, y_1, y_2, \dots, y_t = v \in V_2$. Therefore $y, y_1, y_2, \dots, y_{t-1} \in V_1$. Also $d(y_{t-1}, V_2) = 1$, a contradiction.

(Since $d(z, V_2) = 1$). Suppose V_1 contains two vertices z, x . $d(z, V_2) = 1$. Let $d(x, V_2) = t$. Then $t \geq 2$. There exists a shortest path $x, x_1, x_2, \dots, x_t = v \in V_2$. Since $t \geq 2, x_{t-1} \in V_1$ and $d(x_{t-1}, V_2) = 1$, a contradiction. If $V_1 = \{z\}$ and if $\langle V_2 \rangle$ has an isolate say u , then u and z are adjacent and any path from u to z must contain u . Hence u is not an isolate of $\langle V_2 \rangle$, a contradiction. Therefore $pd_{is}(G) \geq 3$. The

converse is obvious.

Theorem 2.4. $pd_{is}(G) = n$ if and only if $V(G)$ can be partitioned into subsets V_1 and V_2 such that $G = \langle V_1 \rangle + \langle V_2 \rangle$. If any of $\langle V_1 \rangle$ and $\langle V_2 \rangle$ is connected, then its diameter less than or equal to 2 and if one of them is disconnected, then it is totally disconnected.

Proof: Suppose $V(G)$ can be partitioned into subsets V_1 and V_2 such that $G = \langle V_1 \rangle + \langle V_2 \rangle$.

Case(i): Let one of $\langle V_1 \rangle, \langle V_2 \rangle$ be independent and other is connected with diameter less than or equal 2. Let $diam(\langle V_2 \rangle) = 2$. Let $x, y \in V_1$. Let $\pi = \{\{x, y\}, \{x_i\}\}$ where x_i runs over all the vertices of $V_1 - \{x, y\}$ and V_2 . $d(x, x_i) = 2 = d(y, x_i)$ for all $x_i \in V_1 - \{x, y\}$ $d(x, x_i) = 1 = d(y, x_i)$ for all $x_i \in V_2$. Therefore π is not a resolving partition. Suppose $x, y \in V_2$. Then $d(x, x_i) = 1 = d(y, x_i)$ for all $x_i \in V_1$. Let $z \in V_2 - \{x, y\}$. If z is adjacent with y , then there exists a path x, z, \dots, u, y of length greater than or equal to 3, a contradiction. (Since $diam(\langle V_2 \rangle) = 2$). Therefore z is adjacent with both x and y . (If z is not adjacent with both x and y , then $diam(\langle V_2 \rangle) = 2$). Therefore $d(x, z) = d(y, z)$. Therefore π is not a resolving partition. Therefore $pd_{is}(G) = n$. Let $diam(\langle V_2 \rangle) = 1$. Then $\langle V_2 \rangle$ is complete. Therefore no two points of V_2 can be independent. Therefore $pd_{is}(G) = n$.

Case(ii): Let both $\langle V_1 \rangle$ and $\langle V_2 \rangle$ be independent. Then $pd_{is}(G) = n$.

Case(iii): Suppose $\langle V_1 \rangle$ is independent and $\langle V_2 \rangle$ is disconnected but not totally disconnected.

Let $V_1 = \{u_1, u_2, \dots, u_k\}$. Let $V_2 = \{\{u_{k+1}, u_{k+2}, \dots, u_n\}\}$. Let H_1 and H_2 be the components of $\langle V_2 \rangle$. Since $\langle V_2 \rangle$ is not totally disconnected, at least one of H_1, H_2 contains at least two vertices. Let H_1 contain at least two vertices. Since H_1 is connected, there exists two adjacent vertices in H_1 say x, z . Let $y \in V(H_2)$. Let $\pi = \{\{x, y\}, \{x_i\}\}$ where x_i runs over all the vertices of $V_1, H_1 - \{x\}, H_2 - \{y\}$. Since $z \in V(H_1), \{z\} \in \pi$. $d(x, z) = 1, d(y, z) = 2$. Therefore π is an isolate vertex resolving partition of G , a contradiction, since $|\pi| = n - 1$.

Case(iv): Suppose $\langle V_1 \rangle$ and $\langle V_2 \rangle$ are connected and $diam(\langle V_1 \rangle) \leq 2, 1 \leq i \leq 2$. If $diam(\langle V_1 \rangle) = 1$ for $i = 1, 2$ then G is complete and hence $pd_{is}(G) = n$. If $diam(\langle V_1 \rangle) = 2$ or $diam(\langle V_2 \rangle) = 2$ then proceeding as in Case(i), we get that $pd_{is}(G) = n$.

Case(v): Suppose $\langle V_1 \rangle$ is connected and $\langle V_2 \rangle$ is disconnected but not totally disconnected. Then proceeding as in Case (iii), we get a contradiction.

Conversely, suppose $pd_{is}(G) = n$.

Suppose $V(G)$ cannot be partitioned into subsets V_1 and V_2 such that $G = \langle V_1 \rangle + \langle V_2 \rangle$.

Suppose G is complete. Then G can be partitioned into subsets V_1 and V_2 such that $G = \langle V_1 \rangle + \langle V_2 \rangle$.

Suppose G is not complete. Let $\{u_1, u_2, \dots, u_k\}$ be a maximum independent set of G . Then $k \geq 2$.

Let $V_1 = \langle \{u_1, u_2, \dots, u_k\} \rangle$ and $V_2 = \langle V - V_1 \rangle$. Since $G \neq \langle V_1 \rangle + \langle V_2 \rangle$, there exists $u_i \in V_1$ and $y \in V_2$ such that u_i and y are not adjacent. Since u_i is not an isolate of G and since u_i is an isolate of $\langle V_1 \rangle$, u_i is adjacent with some vertex say $z \in V_2$. Then $\pi = \{\{y, z\}, \{x_i\}\}$ where $x_i \in V_1$ or $x_i \in V_2 - \{y, z\}$, is an isolate vertex resolving partition. Since y, z are resolved by u_i . Therefore $pd_{is}(G) \leq n - 1$, a contradiction. Therefore G can be partitioned into subsets V_1 and V_2 such that $G = \langle V_1 \rangle + \langle V_2 \rangle$.

Suppose $\langle V_1 \rangle$ and $\langle V_2 \rangle$ are connected. Suppose $diam(\langle V_1 \rangle) \geq 3$. (similar proof if $diam(\langle V_2 \rangle) \geq 3$). Then there exists a path $u = u_0, u_1, \dots, u_k = v$ in $\langle V_1 \rangle$ where $k = diam(\langle V_1 \rangle) \geq 3$. Let $\pi = \{\{u_i, u_k\}, \{x_i\}\}$ where $x_i \in V_1$

– $\{u_i, u_k\}$, $x_i \in V_2$. u_i, u_k are resolved by u_i . Therefore $pd_{is}(G) \leq n - 1$, a contradiction. Therefore $diam(\langle V_1 \rangle) \leq 2$.

Suppose V_1 is independent and $\langle V_2 \rangle$ is connected. Suppose $diam(\langle V_2 \rangle) \geq 3$. Then proceeding as above $pd_{is}(G) \leq n - 1$, a contradiction. Therefore $diam(\langle V_2 \rangle) \leq 2$. Therefore $G = \langle V_1 \rangle + \langle V_2 \rangle$ and if any of $\langle V_1 \rangle$ and $\langle V_2 \rangle$ is connected, then its diameter less than or equal to 2 and if one of them is disconnected, then it is totally disconnected.

References

- [1]. G. Chartrand, E. Salehi and P. Zhang: The partition dimension of a graph, *Aequationes Math.* 59(2000), 45 -54.
- [2]. G. Chartrand, L. Eroh, M. Johnson and O.R. Oellermann: Resolvability in graphs and the metric dimension of a graph, *Discrete Applied Mathematics*, Vol. 105, Issues 1-3, pp. 99 - 113, (2000).
- [3]. G. Chartrand, F. Okamoto and P. Zhang, The metric chromatic number of a graph, *Australasian Journal of Combinatorics*, 44 (2009), 273- 286.
- [4]. S. Hemalathaa, A. Subramanian, P. Aristotle and V. Swaminathan: Equivalence Resolving Partition of a Graph, communicated.
- [5]. V. Saenpholphat and P. Zhang, Connected partition dimensions of graphs, *Discuss. Math. Graph Theory*, 22(2002), 305-323.
- [6]. P. J. Slater: Leaves of trees, in: *Proc 6th Southeast Conf. Comb., Graph Theory, Comput.*; Boca Raton, 14 (1975), 549 - 559.
- [7]. P. J. Slater: Dominating and reference sets in graphs, *J. Math. Phys. Sci.* 22(1988), 445 - 455.